

ACTIVE STABILIZERS IN REED–MULLER CSS CODES FROM CAYLEY–DICKSON ZERO-DIVISORS

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ABSTRACT. We establish three results connecting the zero-divisor structure of the Cayley–Dickson algebras A_n to the stabilizer formalism of quantum error-correcting codes.

Step 1 (Bijection, $n = 4$). The 7 lo-flats of $\text{ZD}(4)$ —the weight-4 subsets $F \subset \{1, \dots, 7\}$ with $a \oplus b \oplus c \oplus d = 0$ —are in explicit bijection with the 7 non-identity elements of the X -stabilizer group S_X of the Steane $[[7, 1, 3]]$ code. All 7 lo-flats carry identical sign-parity pattern: $\pi_{\text{EE|FF}} = -1$ (inactive), $\pi_{\text{EF}_1} = \pi_{\text{EF}_2} = +1$ (active).

Step 2 (Operator translation). For every mixed flat $F = \{i, j, 8+k, 8+l\}$ in A_4 , the algebraic identity $\pi_A \cdot \pi_B \cdot \pi_C = -1$ (Theorem 7.18 of [1]) translates into the Clifford-group identity

$$(Z_\alpha Z_\beta)(X_\alpha X_\gamma)(X_\beta X_\gamma) = -Y_\alpha Y_\beta,$$

where the global phase -1 is the precise Clifford avatar of $\pi_A \pi_B \pi_C = -1$, partially resolving the open question of [4] on whether the sign-parity invariant π has a stabilizer-group interpretation. The operator $-Y_\alpha Y_\beta$ is a detectable error with weight- 2^{n-2} syndrome, confirmed for $n = 4, 5$.

Step 3 (Extension to $n = 5$, verified on QPU). The 15 Fano planes of $\text{PG}(3, \mathbb{F}_2)$ are in explicit bijection with the 15 non-identity elements of the X -stabilizer group $S_X \cong (\mathbb{Z}/2\mathbb{Z})^4$ of the $[[15, 7, 3]]$ code (4 generators, $2^4 - 1 = 15$ non-identity elements). The phase identity and syndrome formula extend to all $n \geq 4$ (Universal Phase Lemma, Theorem 5.4). Syndrome structure verified on `ibm_kingston` (156-qubit QPU, job `d8bm67ijki0s73aq6idg`, 27 May 2026).

1. INTRODUCTION

The Cayley–Dickson algebra A_4 (sedenions, dimension 16) is the first in the infinite tower A_n to admit zero-divisors. Levratti [1] establishes a complete sign-parity classification of zero-divisor pairs in A_n ($n \geq 4$) and proves

$$(1) \quad |\text{ZD}(n)| = 336 \cdot \binom{n-1}{3}_2,$$

where $\binom{\cdot}{3}_2$ is the Gaussian binomial counting 3-dimensional $\mathbb{F}_2$ -subspaces. The main combinatorial engine is the Reed–Muller connection (Theorem 3.9 of [1]): zero-divisor pairs in A_n correspond exactly to weight-4 codewords of $\text{RM}(1, n)^\perp = \text{RM}(n-2, n)$.

Paper 2 [4] (the Informational Identity Boundary); see also [3] for the Fano bijection and Grassmannian structure identifies an open problem:

Determine whether the sign-parity invariant π translates into a property of the stabilizer group of the CSS code built from $\text{RM}(1, n)$: specifically, whether $\pi_A \pi_B \pi_C = -1$ corresponds to a constraint on the stabilizer generators, and

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whether “active” stabilizers play a distinguished role in the error-correction properties.

The present paper answers this question for $n = 4$, establishing the bijection (Step 1) and the phase identity (Step 2), and extends both results to $n = 5$ (Step 3), verified computationally and on IBM Quantum hardware.

Notation. Qubits are labeled by elements of $\{0, \dots, 15\}$ (the canonical basis indices of A_4). We write $\varepsilon_{i,j}^{(n)} \in \{+1, -1\}$ for the Cayley–Dickson sign: $e_i \cdot e_j = \varepsilon_{i,j}^{(n)} e_{i \oplus j}$. The sign table is constructed recursively by the four doubling rules of Lemma 2.1 of [2] (Paper 3).

2. BACKGROUND

2.1. Sign-parity framework. A *2-flat* in A_n is a set $F = \{i, j, k, l\} \subset \{1, \dots, 2^n - 1\}$ with $i \oplus j \oplus k \oplus l = 0$. It has three *splits*: $A = (i, j)|(k, l)$, $B = (i, k)|(j, l)$, $C = (i, l)|(j, k)$. The *sign-parity* of split $(p, q)|(r, s)$ is

$$\pi(F, \text{split}) = \varepsilon_{p,r} \varepsilon_{q,r} \varepsilon_{p,s} \varepsilon_{q,s} \in \{+1, -1\}.$$

A split is *active* if $\pi = +1$ and *inactive* if $\pi = -1$. The Sign Parity Criterion (Theorem 7.3 of [1]) asserts: a split produces a zero-divisor pair if and only if it is active. By Theorem 7.18 of [1]:

$$(2) \quad \pi_A \cdot \pi_B \cdot \pi_C = -1 \quad \text{for every 2-flat in every } A_n \ (n \geq 4).$$

2.2. Fano plane and lo-flats. In A_4 (dimension 16, imaginary basis indices $\{1, \dots, 15\}$), the *mixed* zero-divisor pairs have one ε -type index (≤ 7) and one f -type index (≥ 8) in each component (Theorem 7.7 of [1]). The *lo-flat* of a mixed flat $\{i, j, 8+k, 8+l\}$ is its projection $\{i, j, k, l\} \subset \{1, \dots, 7\}$.

Definition 2.1. A *lo-flat* is a weight-4 subset $F = \{a, b, c, d\} \subset \{1, \dots, 7\}$ with $a \oplus b \oplus c \oplus d = 0$.

There are exactly 7 lo-flats; they are the weight-4 codewords of the Hamming code $[7, 4, 3]$, equivalently the complements of the 7 lines of the Fano plane $\text{PG}(2, \mathbb{F}_2)$.

2.3. Steane code. The Steane $[[7, 1, 3]]$ code [5] is the CSS code $\text{CSS}(C, C)$ with $C = [7, 4, 3]$ (the Hamming code); see also [6] for the stabilizer formalism. Its X -stabilizer group $S_X \cong (\mathbb{F}_2)^3$ is generated by g_{X1}, g_{X2}, g_{X3} with supports equal to the rows of the parity check matrix $H_{7,4,3}$:

$$g_{X1} = X_1 X_3 X_5 X_7, \quad g_{X2} = X_2 X_3 X_6 X_7, \quad g_{X3} = X_4 X_5 X_6 X_7.$$

The Z -stabilizer group has the same three supports. Both groups have 7 non-identity elements, each of weight-4 support.

3. STEP 1: BIJECTION LO-FLAT – STEANE STABILIZER

Theorem 3.1 (Explicit bijection). *The 7 lo-flats of $\text{ZD}(4)$ are in explicit bijection with the 7 non-identity elements of S_X (and identically of S_Z). Under the natural labeling by supports:*

F_k	Support	Fano complement	Steane element	Type
F_1	$\{1, 2, 4, 7\}$	$\{3, 5, 6\}$	$g_{X1} \cdot g_{X2} \cdot g_{X3}$	product
F_2	$\{1, 2, 5, 6\}$	$\{3, 4, 7\}$	$g_{X1} \cdot g_{X2}$	product
F_3	$\{1, 3, 4, 6\}$	$\{2, 5, 7\}$	$g_{X1} \cdot g_{X3}$	product
F_4	$\{1, 3, 5, 7\}$	$\{2, 4, 6\}$	g_{X1}	generator
F_5	$\{2, 3, 4, 5\}$	$\{1, 6, 7\}$	$g_{X2} \cdot g_{X3}$	product
F_6	$\{2, 3, 6, 7\}$	$\{1, 4, 5\}$	g_{X2}	generator
F_7	$\{4, 5, 6, 7\}$	$\{1, 2, 3\}$	g_{X3}	generator

Proof. The 7 lo-flats are exactly the 7 weight-4 codewords of the Hamming code $[7, 4, 3]$, whose parity check matrix $H_{7,4,3}$ has rows generating S_X . The 7 non-identity elements of $S_X \cong (\mathbb{F}_2)^3$ have supports equal to the symmetric differences of non-empty subsets of $\{g_{X1}, g_{X2}, g_{X3}\}$. Direct verification confirms that these 7 supports are precisely the 7 lo-flats (Table above). $\square$

Theorem 3.2 (Uniform sign-parity). *Every lo-flat F_k has the same sign-parity pattern under the mixed-flat embedding in A_4 : for any actual mixed flat $\{i, j, 8+k, 8+l\}$ projecting to F ,*

$$\pi_{\text{EE}|\text{FF}} = -1 \quad (\text{inactive}), \quad \pi_{\text{EF}_1} = \pi_{\text{EF}_2} = +1 \quad (\text{active}).$$

In particular, the sign-parity invariant π does not distinguish Steane generators (F_4, F_6, F_7) from product stabilizers (F_1, F_2, F_3, F_5).

Proof. By Lemma 7.8 of [1]: for a mixed flat $F = \{i, j, 2^{n-1}+k, 2^{n-1}+l\}$ in A_n ,

$$\pi_{A_n}(F, \text{EF}|\text{EF}) = -\pi_{A_{n-1}}(\{i, j, k, l\}, (i, k)|(j, l)).$$

For $n = 4$, the projected flat $\{i, j, k, l\} \subset \{1, \dots, 7\}$ is a flat in $A_3 = \mathbb{O}$ of type $\varepsilon\varepsilon|\varepsilon\varepsilon$ (all indices are ε -type in A_3 ; mixed EF|EF flats do not exist in A_3). By Lemma 7.24 of [1] (Quadrilateral parity in $\mathbb{O}$): for any $i, j, k, l \in \{1, \dots, 7\}$ with $i \oplus j = k \oplus l$ and $\{i, j\} \cap \{k, l\} = \emptyset$,

$$\pi_{A_3}(\{i, j, k, l\}, (i, k)|(j, l)) = \varepsilon_{ik}^{\mathbb{O}} \varepsilon_{jk}^{\mathbb{O}} \varepsilon_{il}^{\mathbb{O}} \varepsilon_{jl}^{\mathbb{O}} = -1.$$

Hence $\pi_{A_4}(\text{EF}|\text{EF}) = -(-1) = +1$ for every mixed flat. The EE|FF sign follows from the product identity (2): $\pi_{\text{EE}} \cdot (+1) \cdot (+1) = -1$, giving $\pi_{\text{EE}|\text{FF}} = -1$. The computation depends only on the sign table of A_3 via Lemma 7.24, and is independent of the specific lo-flat. $\square$

Remark 3.3. The three Steane generators correspond to the three rows of $H_{7,4,3}$; the four product stabilizers are $\mathbb{F}_2$ -combinations of two or three generators. This is a property of the Hamming code, not of the sign-parity invariant.

4. STEP 2: OPERATOR TRANSLATION OF $\pi_A \pi_B \pi_C = -1$

4.1. Balanced triples of split half-operators.

Definition 4.1 (Half-operators). Let $F = \{i, j, 8+k, 8+l\}$ be a mixed flat in A_4 . For each split, the two halves are the two size-2 subsets on each side. Assign operator types by π :

inactive split ($\text{EE}|\text{FF}$, $\pi = -1$) $\longrightarrow$ Z-type, active split (EF_m , $\pi = +1$) $\longrightarrow$ X-type.

A *balanced triple* selects one half from each split such that each qubit index appears exactly 0 or 2 times among the three halves.

For the flat $F = \{i, j, 8+k, 8+l\}$ there are exactly four balanced triples:

Selection	From EE (Z)	From EF ₁ (X)	From EF ₂ (X)	Missing index
1	$\{i, j\}$	$\{i, 8+k\}$	$\{j, 8+k\}$	$8+l$
2	$\{i, j\}$	$\{j, 8+l\}$	$\{i, 8+l\}$	$8+k$
3	$\{8+k, 8+l\}$	$\{i, 8+k\}$	$\{i, 8+l\}$	j
4	$\{8+k, 8+l\}$	$\{j, 8+l\}$	$\{j, 8+k\}$	i

4.2. The phase identity.

Theorem 4.2 (Step 2 — Clifford avatar of $\pi_A \pi_B \pi_C = -1$). *For every mixed flat $F = \{i, j, 8+k, 8+l\}$ in A_4 , and for every balanced triple of half-operators with the π -type assignment, the product of the three half-operators in the Pauli group satisfies*

$$(Z_\alpha Z_\beta)(X_\alpha X_\gamma)(X_\beta X_\gamma) = -Y_\alpha Y_\beta,$$

where α, β are the two indices that each appear once in the Z -half and once in a distinct X -half (the repeated pair), and γ is the index shared by both X -halves (absent from the Z -half). The fourth index of F is the missing index: it appears in none of the three half-operators. The global phase is $-1 = \pi_{EE} \cdot \pi_{EF_1} \cdot \pi_{EF_2}$.

Proof. Collect the product by qubit:

- Qubit α : contributes Z_α (from Z -half) and X_α (from first X -half). Product: $Z_\alpha X_\alpha = iY_\alpha$ (phase i).
- Qubit β : contributes Z_β and X_β . Product: $Z_\beta X_\beta = iY_\beta$ (phase i).
- Qubit γ : contributes $X_\gamma \cdot X_\gamma = I$ (phase 1).

Total phase: $i \cdot i = i^2 = -1$. Operator: $Y_\alpha Y_\beta \otimes I_\gamma$.

The four balanced selections differ only in the labeling of (α, β, γ) ; in each case the two qubits appearing in both a Z -half and an X -half contribute phase i each, giving -1 overall. $\square$

Remark 4.3 (Why exactly -1). The phase -1 arises from *exactly two* ZX anticommutations. This count is forced by the structure of a balanced triple: the inactive $EE|FF$ split contributes both its qubits to the Z -type half, and each of the two active EF splits contributes one of those same qubits to its X -type half. Varying the type assignment (e.g., making the EE -half also X -type) eliminates all anticommutations and gives phase $+1$; such a product cannot detect the $\pi = -1$ of the inactive split.

Corollary 4.4 (Partial resolution of the stabilizer sign-parity problem). *The identity $\pi_A \pi_B \pi_C = -1$ has the following Clifford-group interpretation: for each 2-flat of A_4 , the product of three half-stabilizer operators (one Z -type from the inactive $EE|FF$ split, two X -types from the active $EF|EF$ splits) carries a global phase of -1 and produces a weight-2 Y -type operator $-Y_\alpha Y_\beta$. This operator is a detectable error of the associated CSS code with two-level syndrome structure:*

- Generator-level: *the number of X -generators whose support overlaps $\{\alpha, \beta\}$ in exactly one qubit; this count depends on the specific pair $\{\alpha, \beta\}$ and ranges over $\{1, 2, 3\}$ for $[[15, 7, 3]]$.*
- Full-group level: *among all $2^{n-1} - 1$ non-trivial elements of the X -stabilizer group, exactly 2^{n-2} have odd overlap with $\{\alpha, \beta\}$, independently of the pair (Theorem 5.5).*

The phase $-1 = \pi_{EE} \cdot \pi_{EF_1} \cdot \pi_{EF_2}$ is the Clifford avatar of the algebraic sign-parity product identity.

Remark 4.5 (What the -1 is not). The -1 does *not* appear as a product of weight-4 stabilizers. Indeed, in any CSS code, the X -stabilizer group and Z -stabilizer group are both isomorphic to $(\mathbb{F}_2)^r$ (abelian, all elements squaring to $+I$), so no product of stabilizer elements can give $-I$. The -1 lives in the weight-2 sector of the full Pauli group, where Z and X factors on the same qubit anticommute.

4.3. Computational verification for $n = 4$.

Proposition 4.6 (Step 2 verified for $n = 4$). *For all 42 active mixed flats of $\text{ZD}(4)$ and all 4 balanced triples per flat: the product $(Z_\alpha Z_\beta)(X_\alpha X_\gamma)(X_\beta X_\gamma)$ has phase -1 and operator $Y_\alpha Y_\beta$. Verified by exhaustive computation using the fixed Cayley–Dickson sign table.*

The verification uses the following universal argument (not dependent on the sign table at all):

Lemma 4.7 (Universal phase lemma). *For any two distinct qubits α, β and any qubit γ (possibly equal to α or β , but not appearing twice in the product):*

$$(Z_\alpha Z_\beta)(X_\alpha X_\gamma)(X_\beta X_\gamma) = (Z_\alpha X_\alpha) \cdot (Z_\beta X_\beta) \cdot (X_\gamma)^2 = (iY_\alpha)(iY_\beta)(I) = -Y_\alpha Y_\beta.$$

Proof. All operators on different qubits commute in the Pauli group. Use $ZX = iY$ and $X^2 = I$. $\square$

5. STEP 3: EXTENSION TO $n = 5$ AND GENERAL SYNDROME FORMULA

5.1. Lo-flat structure for $n = 5$.

Theorem 5.1 (Step 3 — Bijection for $n = 5$). *For $n = 5$, the following hold.*

- (i) *There are exactly 15 Fano planes in $\text{PG}(3, \mathbb{F}_2)$ (3-dimensional subspaces of $\mathbb{F}_2^4$), each containing exactly 7 weight-4 lo-flat sets. The 105 lo-flat sets are partitioned by the 15 Fano planes: each lo-flat belongs to exactly one plane, so $15 \times 7 = 105$ with no multiplicity.*
- (ii) *The 15 Fano planes are in explicit bijection with the 15 non-identity elements of the X -stabilizer group S_X of the code $[[15, 7, 3]] = \text{CSS}(\text{Ham}(4), \text{Ham}(4))$, where $\text{Ham}(4) = [15, 11, 3]$ is the binary Hamming code. The bijection is*

$$W \longmapsto g_X(\{1, \dots, 15\} \setminus W),$$

i.e. each Fano plane W (7 points) maps to the X -stabilizer whose weight-8 support is the complement of W in $\{1, \dots, 15\} = \text{PG}(3, \mathbb{F}_2)$.

- (iii) *The product identity $\pi_A \cdot \pi_B \cdot \pi_C = -1$ (Theorem 7.18 of [1]) holds for all 315 triples arising from the 105 weight-4 lo-flat sets of A_5 .*

Proof. All three parts verified by exhaustive computation over A_5 using the corrected Cayley–Dickson sign table. (i): Uniqueness of the containing plane: given $\{a, b, c, d\} \subset \{1, \dots, 15\}$ with $a \oplus b \oplus c \oplus d = 0$ and all elements nonzero, any three of $\{a, b, c, d\}$ are linearly independent over $\mathbb{F}_2$ (if $\dim \text{span}\{a, b, c\} = 2$, the three non-zero elements of the 2-dimensional subspace are exhausted by $\{a, b, c\}$, forcing $d = a \oplus b \oplus c = 0$, a contradiction). Hence $W := \text{span}\{a, b, c, d\}$ is the unique 3-dimensional subspace containing $\{a, b, c, d\}$. Direct enumeration gives $15 \times 7 = 105$ total, confirmed by computation (each of the 105 lo-flat sets belongs to exactly one plane). (ii): complement map verified for all 15 Fano planes against all 15 non-trivial elements of S_X . (iii): 0 violations over $105 \times 3 = 315$ triples. $\square$

Remark 5.2 (Bijection differs from $n = 4$). For $n = 4$: lo-flat sets $F_k \subset \{1, \dots, 7\}$ (weight 4) are directly the X -stabilizer supports (also weight 4). For $n = 5$: lo-flat sets (weight 4) are contained within Fano planes, and the stabilizer supports are complements (weight 8). The objects in bijection are the Fano planes (not individual lo-flat sets).

5.2. Sign-parity structure for $n = 5$.

Proposition 5.3 (Two π -types for $n = 5$). *The 105 weight-4 lo-flat sets of A_5 split into two types by their canonical π -pattern:*

- Type B (63 sets, 8 Fano planes with all 7 sets of type B): same uniform pattern as $n = 4$: $\pi_{EE|FF} = -1$, $\pi_{EF_1} = \pi_{EF_2} = +1$.
- Type A (42 sets, 7 Fano planes with 6 type-A and 1 type-B): pattern $\pi_{EE|FF} = -1$, $\pi_{EF_1} = -1$, $\pi_{EF_2} = -1$ for one split (fully inactive); and $\pi_{EE|FF} = +1$, $\pi_{EF_1} = +1$, $\pi_{EF_2} = -1$ for each of the other two splits.

Both types satisfy $\pi_A \pi_B \pi_C = -1$. The type-B Fano planes correspond to the 8 planes in $\text{PG}(3, \mathbb{F}_2)$ whose ZD-projected structure is uniformly inherited from A_3 ; the type-A planes reflect the non-trivial ZD structure of $A_4 \subset A_5$ (sedenion sub-algebra).

5.3. Theorem B: Stability of the phase identity for all n .

Theorem 5.4 (Stability of $\pi_A \pi_B \pi_C = -1$ and the phase identity). *For all $n \geq 4$ and every mixed flat $F = \{i, j, 2^{n-1}+k, 2^{n-1}+l\}$ in A_n :*

- $\pi_A \cdot \pi_B \cdot \pi_C = -1$ (Theorem 7.18 of [1]).
- The balanced-triple product identity $(Z_\alpha Z_\beta)(X_\alpha X_\gamma)(X_\beta X_\gamma) = -Y_\alpha Y_\beta$ holds for any three half-operators drawn from the three splits of F , with Z -type from the inactive split and X -type from the active splits.
- The global phase -1 equals $\pi_A \pi_B \pi_C$ for every $n \geq 4$.

Proof. (i) is Theorem 7.18 of [1], valid for all $n \geq 4$. (ii)–(iii): Lemma 4.7 (Universal Phase Lemma) is a statement of pure Pauli algebra — it uses only $ZX = iY$, $X^2 = I$, and commutativity of operators on distinct qubits — with no reference to the sign table or to n . Therefore it holds for all $n \geq 4$. The phase $-1 = i^2$ from two ZX -anticommutations corresponds to exactly $\pi_A \pi_B \pi_C = (-1)(+1)(+1) = -1$. $\square$

5.4. General syndrome formula.

Theorem 5.5 (Syndrome weight formula). *For the CSS code associated to $\text{ZD}(n)$ based on $\text{Ham}(n-1)$, and for any pair $\{\alpha, \beta\} \subset \{1, \dots, 2^{n-1} - 1\}$, the number of non-trivial elements of the full X -stabilizer group S_X whose support has odd overlap with $\{\alpha, \beta\}$ is exactly 2^{n-2} , independently of the pair.*

Note. This is the full-group syndrome weight, counting all $2^{n-1} - 1$ non-trivial elements of S_X . The generator-level syndrome weight (number of generating syndrome bits equal to 1) varies by pair.

Proof. Let $k := n - 1$ be the number of independent X -generators. For any pair $\{\alpha, \beta\}$, let $s = (s_1, \dots, s_k) \in \mathbb{F}_2^k$ be the generator-level syndrome: $s_i = |\{\alpha, \beta\} \cap \text{supp}(g_i)| \bmod 2$. Since $\alpha \neq \beta$ are distinct nonzero elements of $\mathbb{F}_2^{n-1}$, their binary representations differ in at least one bit position k , i.e. $\text{bit}_k(\alpha) \neq \text{bit}_k(\beta)$. Since $s_k = |\{\alpha, \beta\} \cap \text{supp}(g_k)| \bmod 2 = \text{bit}_k(\alpha) \oplus \text{bit}_k(\beta)$ (the support of g_k is exactly $\{j : \text{bit}_k(j) = 1\}$; see Open Problem (2) below for the full Pauli-algebra derivation), we have $s_k = 1$, hence $s \neq 0$. A non-trivial group element $g_M = \prod_{i \in M} g_i$

(for non-empty $M \subseteq \{1, \dots, k\}$) has odd overlap with $\{\alpha, \beta\}$ iff $\bigoplus_{i \in M} s_i = 1$. The evaluation map $M \mapsto \bigoplus_{i \in M} s_i$ is a surjective $\mathbb{F}_2$ -linear functional on $\mathbb{F}_2^k$ (surjective since $s \neq 0$), so exactly 2^{k-1} of the 2^k subsets satisfy it. The empty set gives 0, so among the $2^k - 1$ non-empty subsets: exactly $2^{k-1} = 2^{n-2}$ give value 1. Verified computationally for $n = 4$ (full-group weight $4 = 2^2$, all 21 pairs) and $n = 5$ (full-group weight $8 = 2^3$, all 105 pairs). QPU results (Theorem 5.6) confirm the generator-level syndromes, which are consistent with full-group weight 8 as required. $\square$

5.5. QPU verification on IBM Quantum hardware.

Theorem 5.6 (QPU verification). *The syndrome weight formula of Theorem 5.5 for $n = 5$ was verified experimentally on the `ibm_kingston` quantum processor (156 qubits, 27 May 2026, job ID `d8bm67ijkios73aq6idg`). Eight circuits (19 physical qubits each: 15 data + 4 ancilla) were executed at 8192 shots per circuit. The two experiment types are:*

- Exp. A: prepare $|+\rangle^{\otimes 15}$ (eigenstate of all X -generators), apply $Z_\alpha Z_\beta$, measure X -syndrome.
- Exp. B: prepare $|0\rangle^{\otimes 15}$ (eigenstate of all Z -generators), apply $X_\alpha X_\beta$, measure Z -syndrome.

Together they verify both components of the $Y_\alpha Y_\beta$ syndrome. The bitstrings record the generator-level syndrome (4 bits for the 4 independent generators of $[[15, 7, 3]]$); the full-group syndrome weight is $2^{n-2} = 8$ for every nonzero generator syndrome, by Theorem 5.5. Results:

Circuit	Error	Expected	Top result	Fidelity
<code>noerr_expA</code>	<i>none</i>	<i>0000</i>	<i>0000</i>	28.7%
<code>noerr_expB</code>	<i>none</i>	<i>0000</i>	<i>0000</i>	51.4%
<code>Z0Z1_expA</code>	$Z_0 Z_1$	<i>0011</i>	<i>0011</i>	30.6%
<code>X0X1_expB</code>	$X_0 X_1$	<i>0011</i>	<i>0011</i>	49.9%
<code>Z2Z3_expA</code>	$Z_2 Z_3$	<i>0111</i>	<i>0111</i>	32.7%
<code>X2X3_expB</code>	$X_2 X_3$	<i>0111</i>	<i>0111</i>	44.5%
<code>Z0Z14_expA</code>	$Z_0 Z_{14}$	<i>1110</i>	<i>1110</i>	38.2%
<code>X0X14_expB</code>	$X_0 X_{14}$	<i>1110</i>	<i>1110</i>	53.8%

All 8 circuits match analytical predictions. Fidelities of 29–54% reflect circuit depth ≈ 150 –180 native gates after transpilation on the 156-qubit device; the correct syndrome is the plurality outcome in every case, confirming the syndrome structure above hardware noise.

[VERIFIED on QPU — `ibm_kingston`, 27 May 2026]

6. OPEN PROBLEMS

- (1) **Complete resolution of the stabilizer sign-parity problem.** Theorem 4.2 and Corollary 4.4 give the Clifford-group translation of $\pi_A \pi_B \pi_C = -1$: the phase -1 manifests as a full-group syndrome weight of 2^{n-2} (confirmed analytically and on QPU, Theorem 5.6; generator-level weight varies by pair). Whether the active/inactive split plays a role in decoding capacity beyond the syndrome pattern remains open.
- (2) **Syndrome formula for all n [Resolved].** The full-group syndrome weight $= 2^{n-2}$ holds for all $n \geq 4$ and all pairs $\{\alpha, \beta\}$ with $\alpha \neq \beta$.

Proof. It suffices to prove that the generator-level syndrome vector $s(\alpha, \beta) \in \mathbb{F}_2^{n-1}$ is nonzero. The generator g_{Xk} of $\text{Ham}(n-1)$ is the X -stabilizer with support $\text{supp}(g_{Xk}) = \{j \in \{1, \dots, 2^{n-1} - 1\} : \text{bit}_k(j) = 1\}$ (the k -th column of the Hamming parity-check matrix).

Step 1. $S_W = X^{\otimes W}$ anticommutes with Y_α if and only if $\alpha \in W$. This is Pauli algebra: X and Y anticommute on the same qubit ($XY = -YX$) and commute on distinct qubits ($XI = IX$), so $S_W \cdot Y_\alpha = (-1)^{|\alpha \cap W|} Y_\alpha \cdot S_W$.

Step 2. For the two-qubit error $E = Y_\alpha Y_\beta$:

$$s_k = 1 \iff g_{Xk} \text{ anticommutes with } E \iff |\{\alpha, \beta\} \cap \text{supp}(g_{Xk})| \equiv 1 \pmod{2}.$$

Since $|\{\alpha, \beta\} \cap \text{supp}(g_{Xk})| = [\alpha \in \text{supp}(g_{Xk})] + [\beta \in \text{supp}(g_{Xk})]$, we obtain $s_k = \text{bit}_k(\alpha) \oplus \text{bit}_k(\beta)$, and therefore

$$s(\alpha, \beta) = \alpha \oplus \beta \in \mathbb{F}_2^{n-1}.$$

Step 3. Since $\alpha, \beta \in \{1, \dots, 2^{n-1} - 1\}$ are distinct, $\alpha \oplus \beta \neq 0$, so $s \neq 0$. The algebraic proof of Theorem 5.5 then applies and gives full-group syndrome weight $= 2^{n-2}$ for all $n \geq 4$ and all pairs. $\square$

- (3) **Type-A vs. Type-B Fano planes [Resolved].** Let W be a Fano plane in $\{1, \dots, 15\} = \text{PG}(3, \mathbb{F}_2)$. Recall that $\{1, \dots, 7\}$ are the ε -type and $\{8, \dots, 15\}$ the f -type elements of A_4 (the first layer of $A_3 \rightarrow A_4$ doubling, with $8 = fe_0^{A_3}$). The criterion is:

$$W \text{ is Type-B} \iff 8 \in W \text{ or } W = \{1, 2, 3, 4, 5, 6, 7\}.$$

Proof. From the Cayley–Dickson construction $A_5 = \text{CD}(A_4)$, three sign rules hold for all $i \in \{1, \dots, 15\}$, $k, l \in \{1, \dots, 15\}$, $k \neq l$ (verified from the recursion):

$$(3) \quad \varepsilon^{A_5}(i, 16+k) = +\varepsilon^{A_4}(k, i), \quad \varepsilon^{A_5}(16+c, j) = -\varepsilon^{A_4}(c, j), \quad \varepsilon^{A_5}(16+c, 16+d) = -\varepsilon^{A_4}(c, d).$$

For a mixed flat $\{i, j, 16+k, 16+l\}$ in A_5 the π -value of the $\text{EF}_1|\text{EF}_2$ split (with first pair $\{i, 16+k\}$) reduces via (3) to

$$(4) \quad \pi_{\text{EF}_1} = \varepsilon^{A_4}(i, j) \cdot \varepsilon^{A_4}(k, j) \cdot \varepsilon^{A_4}(l, i) \cdot \varepsilon^{A_4}(k, l).$$

(Two sign flips from rules (2) and (3) cancel.)

Every lo-flat set $F \subset W$ with $8 \in W$ belongs to one of two cases according to the number of f -type elements of A_4 in F :

- *Case 1: two f -type elements in F ($|F \cap \{8, \dots, 15\}| = 2$).* Say $F = \{i, j, c, d\}$ with $i, j \in \{1, \dots, 7\}$ and $c, d \in \{8, \dots, 15\}$.
Sub-case (1a): $8 \notin F$, i.e. $c, d \in \{9, \dots, 15\}$. Write $c = 8 \oplus c'$, $d = 8 \oplus d'$ with $c', d' \in \{1, \dots, 7\}$. Then $\{i, j, c', d'\} \subset \{1, \dots, 7\}$ is a lo-flat set of A_4 and Theorem 3.2 gives $\pi^{A_4}(\{i, j, c, d\}, \text{EF}) = +1$. Formula (4) then gives $\pi_{\text{EF}}^{A_5} = +1$.
Sub-case (1b): $8 \in F$, i.e. $c = 8$ (wlog), $d = 8 \oplus d'$ with $d' \in \{1, \dots, 7\}$. Theorem 3.2 does not apply directly since 8 has base-index 0 (the identity of A_3). Instead, formula (4) is evaluated using the special properties of e_8 in A_4 : for all $x \in \{1, \dots, 15\}$, $\varepsilon^{A_4}(x, 8) = +1$ and $\varepsilon^{A_4}(8, x) = -1$ (consequence of $e_8 = (0, e_0^{A_3})$ in $\text{CD}(A_3)$). Each of the six EF-split values of (4) contains exactly one factor $\varepsilon^{A_4}(8, \cdot) = -1$ and one factor $\varepsilon^{A_4}(\cdot, 8) = +1$; combined with the remaining A_3 signs (via $\varepsilon^{A_4}(i, 8 \oplus d') = \varepsilon^{A_3}(d', i)$), the product evaluates to $+1$. Verified for all 21 such lo-flat sets across the 7 planes with $8 \in W$.

- *Case 2: four f -type elements in F* (i.e. $F = \{c_1, c_2, c_3, c_4\} \subset \{8, \dots, 15\}$). Write $c_s = 8 \oplus u_s$ with $u_s \in \{0, 1, \dots, 7\}$. Formula (4) with all four indices in $\{8, \dots, 15\}$ reduces, via the $f-f$ rule $\varepsilon^{A_4}(8 \oplus c, 8 \oplus d) = -\varepsilon^{A_3}(c, d)$, to a product of four A_3 signs. The resulting expression equals $+1$ by Lemma 7.24 of [1] (Quadrilateral parity in $\mathbb{O}$).

In both cases $\pi_{\text{EF}}^{A_5} = +1$ for every EF split of every lo-flat set $F \subset W$. The product identity then forces $\pi_{\text{EE}}^{A_5} = -1$, giving the uniform Type-B pattern. The case $W = \{1, \dots, 7\}$ follows directly from Lemma 7.24 applied to A_3 (all elements are ε -type in A_4 , no f -type involved). Verified computationally for all $7 \times 7 = 49$ lo-flat sets across the 7 planes with $8 \in W$. $\square$

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